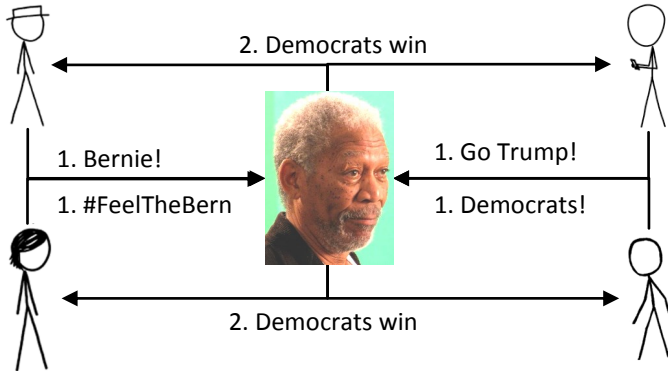


# Efficient Multiparty Computation with Identifiable Abort

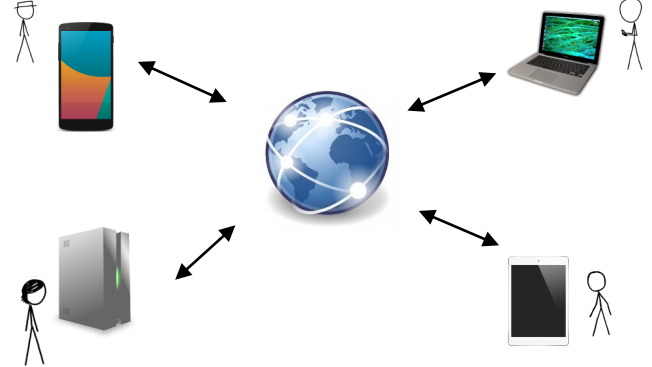
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## What is Multiparty Computation? An Example: Voting

### Dream Election

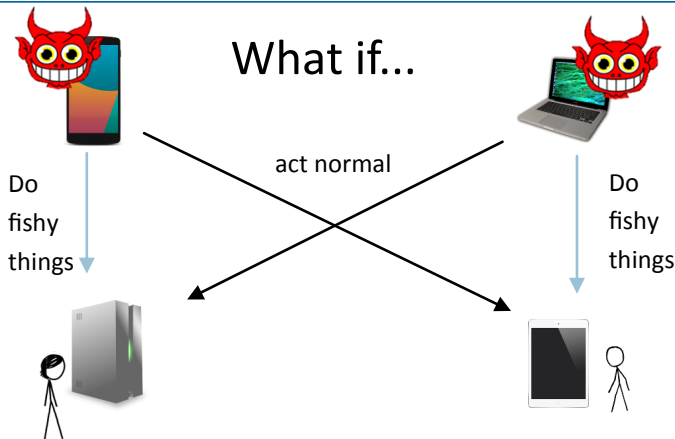


### Real World



**Goal:** Compute a function on inputs from different parties by sending messages over a network. Do this so that the inputs remain private, even if someone arbitrarily deviates from the protocol.

### What if...



- both and disagree on who is a bad guy
- even worse: cannot just trust what says because she may also be a corrupt party!

- the problem is also called **MPC with Identifiable Abort**: If something goes wrong, then we want to identify a cheater

Current solutions to the problem:

- [1] gives a theoretical solution where every party proves in every step that it behaved honestly (inefficient)
- [2] just assumes to abort in this case (this opens up the protocol to denial of service attacks)

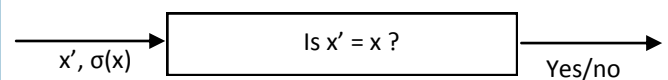
## Our Key Idea

- for each value  $x$  that has, knows a value  $\sigma(x)$  and has a  $\tau(x)$ . Neither nor knows  $x$  though
- when computes  $f(x)$  from  $x$ , then and can compute  $\sigma(f(x))$ ,  $\tau(f(x))$  respectively

Can I add  $x$  and  $y$ ??

Yes! Because can locally obtain a  $\sigma(x+y)$  from  $\sigma(x)$ ,  $\sigma(y)$ ! can also compute  $\tau(x+y)$

has a magic box: When he gets a value  $x'$  from he can check that:



Everyone has this box and all get the same answer! We show how to do this very fast and solve the above problem.

### Bibliography:

- [1] Yuval Ishai, Rafail Ostrovsky & Vassilis Zikas: "Secure Multiparty Computation with Identifiable Abort", CRYPTO 2014
- [2] Ivan Damgård, Valerio Pastro, Nigel Smart and Sarah Zakarias: "Multiparty computation from somewhat homomorphic encryption", CRYPTO 2012
- [3] Carsten Baum, Emmanuela Orsini and Peter Scholl: "Efficient Secure Multiparty Computation with Identifiable Abort", to appear